

# Simulation of random variables

Jean-Michel Marin

University of Montpellier  
Faculty of Sciences

HAX918X / 2024-2025

- 1 Methods involving the uniform distribution on  $[0, 1]$
- 2 The accept-reject algorithm

## Methods involving the uniform distribution on $[0, 1]$

**Proposition** Let  $X$  be a real random variable ( $X(\Omega) \subseteq \mathbb{R}$ ), with cumulative distribution function  $F(x) = \mathbb{P}(X \leq x) = \int_{-\infty}^x f(u) d\mu(u)$

- ▶ If  $F(x)$  is continuous, then  $U = F(X)$  is distributed according to a uniform  $[0, 1]$  distribution
- ▶ Even if  $F(x)$  is not continuous, the inequality  $\mathbb{P}(F(X) \leq t) \leq t$  is true for all  $t \in [0, 1]$
- ▶ If  $F^{[-1]}(y) = \inf\{x : F(x) \geq y\}$  ( $0 < y < 1$ ) and if  $U$  is distributed from a uniform  $[0, 1]$  distribution, then  $F^{[-1]}(U)$  is distributed according to  $F(x)$

# Methods involving the uniform distribution on $[0, 1]$

To perform probabilistic simulations on a computer, a pseudo-random number generator is used

Such a generator returns a sequence  $(x_n)_{n \in \mathbb{N}}$  of real numbers between 0 and 1

These numbers are calculated by a deterministic algorithm but imitate a realization of a sequence of iid uniform  $[0, 1]$  random variables

The good behavior of the sequence is verified by means of statistical tests

## Methods involving the uniform distribution on $[0, 1]$

A standard method to construct the sequence  $(x_n)_{n \in \mathbb{N}}$  is the congruence  $x_n = y_n/N$  where the  $y_n$  are integers such that

$$y_{n+1} = (ay_n + b) \pmod{N}$$

The period of the congruence generator is always smallest than  $N - 1$

The choice of  $a$ ,  $b$  et  $N$  is done such that

- ▶ the period of the generator is the largest as possible
- ▶ the sequence  $(x_n)_{n \in \mathbb{N}}$  is as close as possible to an iid uniform  $[0, 1]$  sequence

# Methods involving the uniform distribution on $[0, 1]$

**Proposition** If  $U \sim \mathcal{U}([0, 1])$  then  $X = a + (b - a)U \sim \mathcal{U}([a, b])$

**Proposition** If  $U \sim \mathcal{U}([0, 1])$  then  $X = \mathbb{I}_{U \leq p} \sim \mathcal{B}(1, p)$

**Proposition** If  $U_1, \dots, U_n$  are  $n$  iid uniform random variables on  $[0, 1]$ , then  $X = \sum_{i=1}^n \mathbb{I}_{U_i \leq p} \sim \mathcal{B}(n, p)$

It is always possible to obtain a simulation following a random variable which takes the values  $(x_i)_{i \in \mathbb{N}^*}$  with respective probabilities  $(p_i)_{i \in \mathbb{N}^*}$  (with  $p_i \geq 0$  such as  $\sum_{i \in \mathbb{N}^*} p_i = 1$ ) using a single uniform variable on  $[0, 1]$

# Methods involving the uniform distribution on $[0, 1]$

**Proposition** If  $U \sim \mathcal{U}([0, 1])$ , then

$$X = x_1 \mathbb{I}_{U \leq p_1} + x_2 \mathbb{I}_{p_1 < U \leq p_1 + p_2} + \dots + x_i \mathbb{I}_{\sum_{j=1}^{i-1} p_j < U \leq \sum_{j=1}^i p_j} + \dots$$

is distributed as a random variable that takes values  $(x_i)_{i \in \mathbb{N}^*}$  with associated probabilities  $(p_i)_{i \in \mathbb{N}^*}$

Requires coding a loop on  $i$  with stopping rule

$\sum_{j=1}^{i-1} p_j < U \leq \sum_{j=1}^i p_j \implies$  it can take a while if the sequence  $(p_i)_{i \in \mathbb{N}^*}$  converges slowly to 1.

# Methods involving the uniform distribution on $[0, 1]$

**Proposition** If  $U_1$  and  $U_2$  are two  $\mathcal{U}([0, 1])$  independent random variables, then

$$X_1 = \sqrt{-2 \ln(U_1)} \cos(2\pi U_2)$$

and

$$X_2 = \sqrt{-2 \ln(U_1)} \sin(2\pi U_2)$$

are two independent standard Gaussian random variables

Recall that if  $X \sim \mathcal{N}(0, 1)$  then  $\mu + \sigma X \sim \mathcal{N}(\mu, \sigma^2)$



# The accept-reject algorithm

**Target** distribution with pdf  $p$  on  $\mathbb{R}^d$

**Instrumental** distribution with pdf  $q$  on  $\mathbb{R}^d$

There exists  $k \geq 1$  such that

$$\forall x \in \mathbb{R}^d, \quad p(x) \leq kq(x)$$

# The accept-reject algorithm

- 0) Set  $i=1$
- 1) Generate  $Y_i$  from  $q$
- 2) Calculate  $M = \frac{p(Y_i)}{kq(Y_i)}$
- 3) Generate  $U_i \sim \mathcal{U}([0, 1])$
- 4) If  $U_i > M$ , then  $i = i + 1$  and back 1)  
If  $U_i \leq M$ , then  $X = Y_i$

# The accept-reject algorithm

Note  $N = \inf\{i \geq 1 : kq(Y_i)U_i \leq p(Y_i)\}$  ( $N$  is a random variable), we have  $X = Y_N$

**Proposition**  $N$  is distributed according to a geometric distribution with parameter  $1/k$ ,  $\mathbb{E}(N) = k$

$N$  is independent of  $(Y_N, kq(Y_N)U_N)$  which is uniformly distributed on

$$\{(x, z) \in \mathbb{R}^d \times \mathbb{R} : 0 \leq z \leq p(x)\}$$

Typically,  $X = Y_N$  is distributed from  $p$